

Cellular Automata as Models of Complexity

Stephen Wolfram, 1984

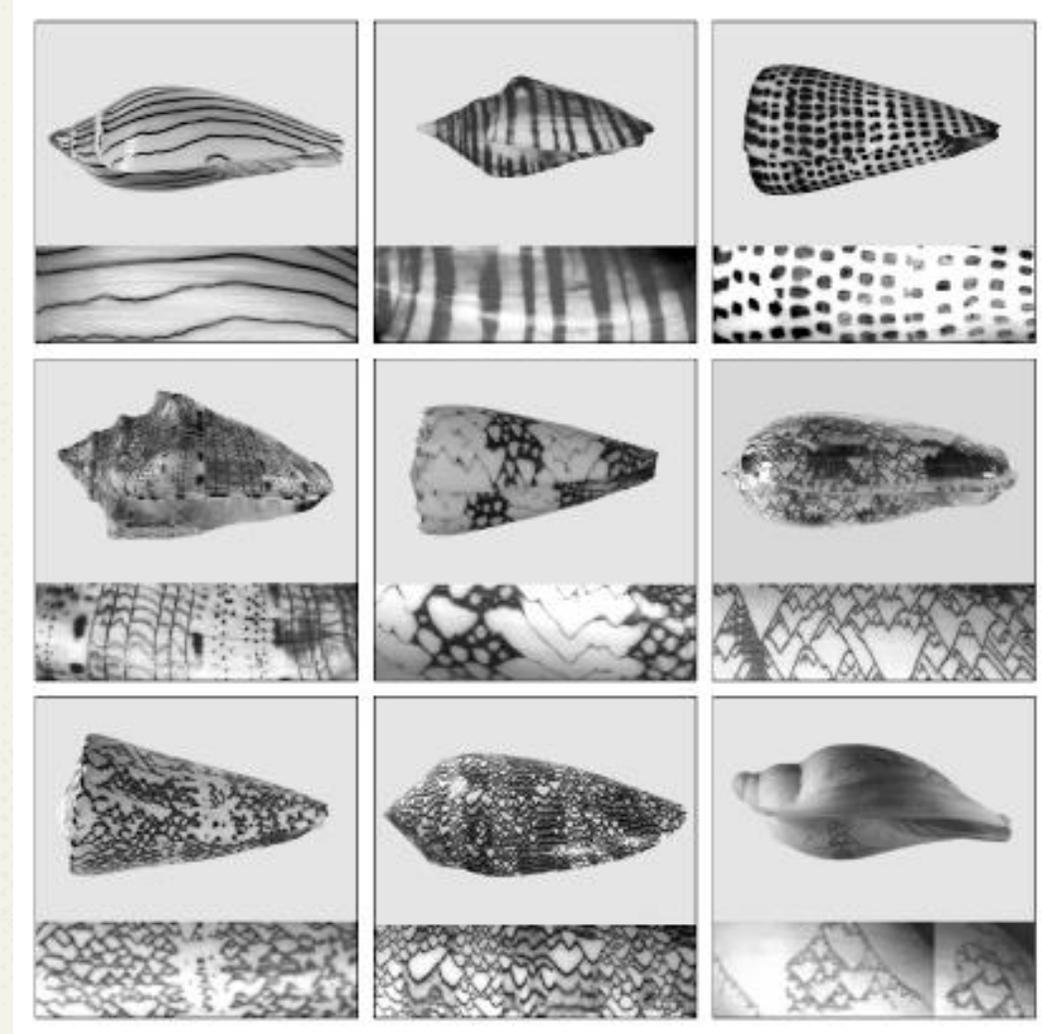
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March 3rd 2017



Introduction

- Cellular Automata
- Applications
- Types of Cellular Automata
- Irreversibility
- Limiting Entropy
- Information Propagation
- Formal Language Theory
- Complexity
- Computation Theory
- Thermodynamics



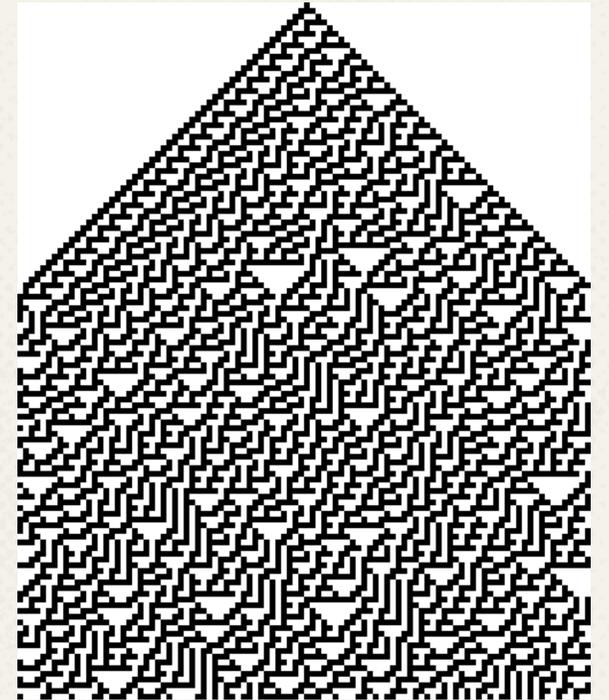
Cellular Automata

A mathematical system constructed from many identical components.

A grid of binary values (or values $0, \dots, k-1$) that update according to a deterministic rule depending on a neighborhood

Complex behavior can arise from simple rules

$$a_i^{(t+1)} = \phi[a_{i-r}^{(t)}, a_{i-r+1}^{(t)}, \dots, a_{i+r}^{(t)}]$$



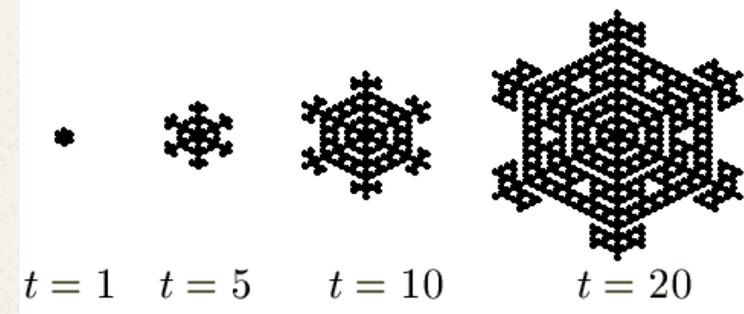
<https://www.cs.rit.edu/~ark/calarge.png>

Applications

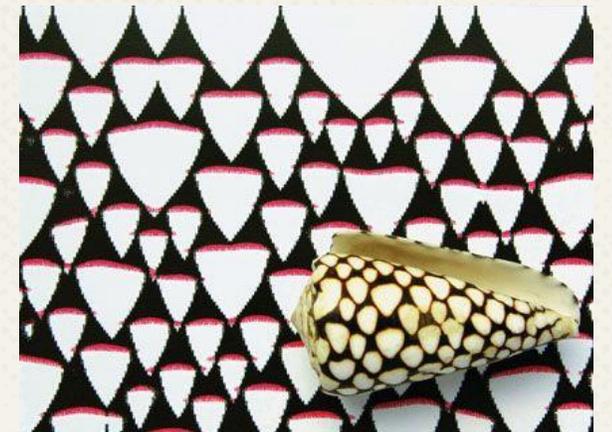
Cellular automata may serve as suitable models for a wide range of natural systems such as:

- The growth of dendritic crystals (snowflakes)
- Pigmentation patterns on mollusk shells

This paper concentrates on general mathematical features of behavior – not specific applications



http://ca.olin.edu/2005/cellular_automata/crystalization_example.gif

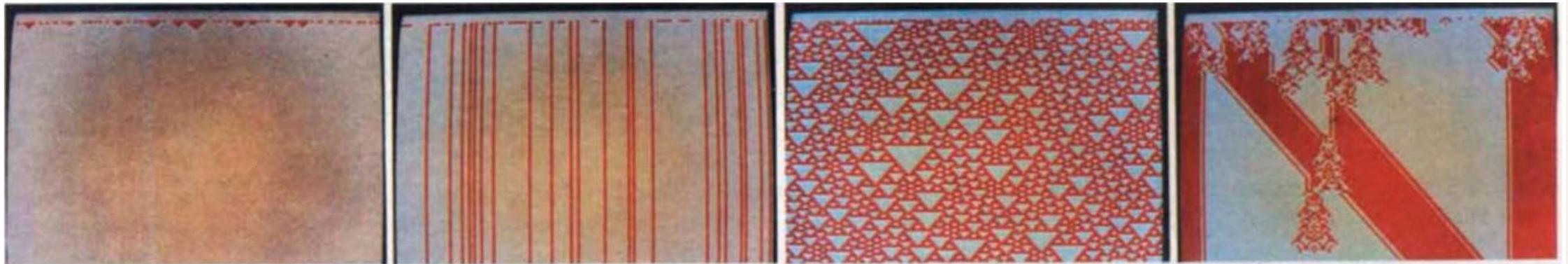


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Cellular Automata

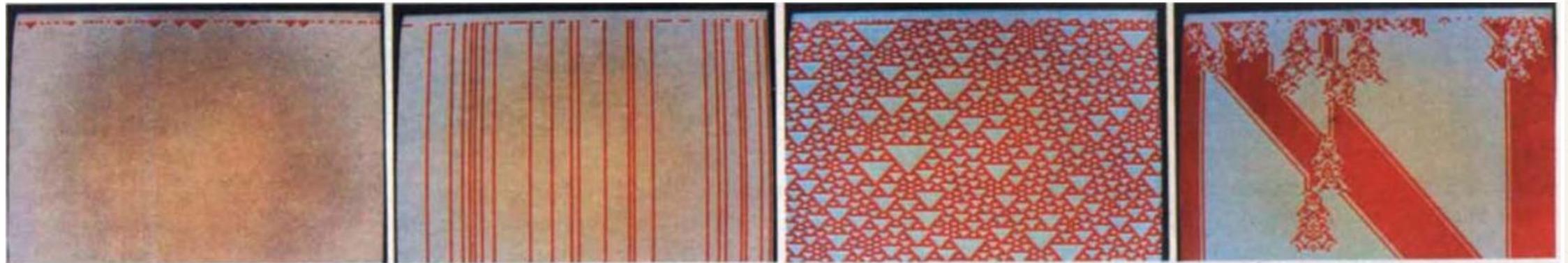
Four classes of automata:

1. Disappears with time – Fixed homogenous state
2. Evolves to a fixed finite size – Limiting cycles
3. Grows indefinitely at a fixed speed – chaotic, aperiodic pattern
4. Grows and contracts irregularly – Complex localized structures



Irreversibility

Different configurations may evolve to the same configuration, thus initial configurations cannot be recalculated



Limiting Entropy (Dimension)

A generalized Measure of the density of configurations generated

$$\text{Limiting Entropy} = d^{(x)} = \lim_{X \rightarrow \infty} \frac{1}{X} \log_k N(X)$$

$N(X)$ gives the number of distinct sequences of X site values that appear

$d(x) = 0$ for class 1 and 2 due to them reaching a limit set

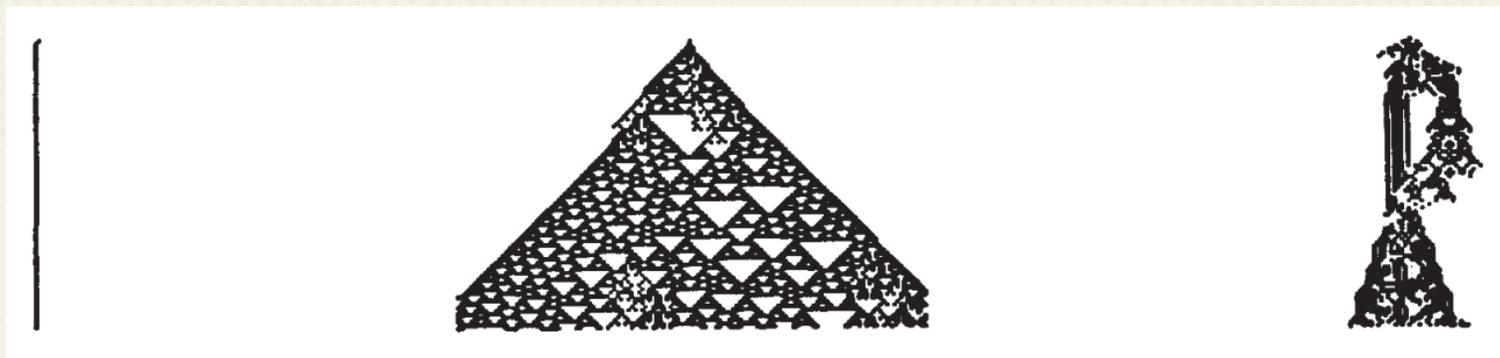
For class 3 the limiting entropy decreases in time, suggesting a fractal subset of all possible configurations occurs

Information Propagation

Characterized by the stability of predictability of their behavior under small changes in initial conditions

Information propagation rates:

- Finite distance in class 1 and 2
- Infinite distance at fixed speed in class 3
- Irregular propagation over an infinite range in class 4



Information Propagation

- The degree of sensitivity to initial conditions shows us degrees of predictability for the cellular automation:
- Class 1 – Entirely predictable, independent of initial state
- Class 2 – Local behavior predictable from local initial state
- Class 3 – Behavior depends on an ever-increasing initial region
- Class 4 – behavior effectively unpredictable

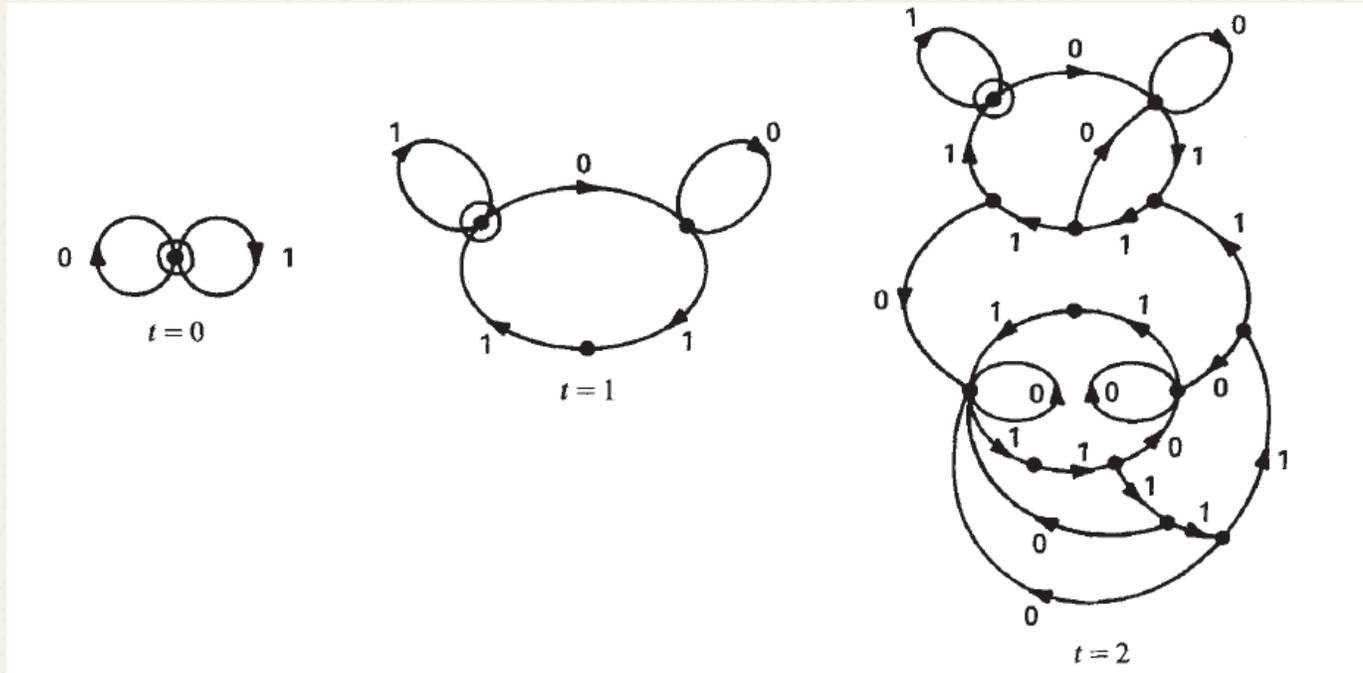
Formal Language Theory

- What is formal language?
 - Set of symbols/strings over some finite alphabet + set of rules (grammar).
 - Recognized by idealized computers with a CPU containing a finite number of internal states + memory
- Types of formal languages:
 - Regular languages (computation model -> Finite State Automaton): No memory required
 - Context-free languages (computation model -> Non-deterministic Pushdown Automaton): Need LIFO stack
 - Context-sensitive languages: memory size = input word size
 - Unrestricted languages (Turning Machine): Large memory required

Why Formal Language Theory?

- Quantities like entropy and dimension (Information theory) give rough characterization of cellular automaton behavior.
- Computation theory gives complete description of self-organization in cellular automaton
- Sets of cellular automaton configurations = Formal Languages
 - Consists of symbols (site values) + set of rules (grammar)
 - Set of all possible initial configurations = Trivial formal language
 - Set of configurations after any finite number of time steps = Regular language

Formal Language Theory



- Set of configurations generated in the first few time steps of evolution according to a class 3 cellular automaton rule.

• Observations from graph •

- Possible configurations = possible paths, beginning at the encircled node.
- All possible configurations are allowed at $t=0$. As time grows, less configurations are generated.
- Set of all possible configurations at each time step forms regular language.
- Number of nodes in the smallest graph = Regular language complexity of the set of configurations.
 - Larger the number of nodes = more complicated sets. E.g. $t=3$, complexity=107 and $t=4$, complexity = 2867

Complexity

- Regular language complexity for the sets generated by cellular automaton evolution is non-decreasing with time.
 - Complexity represents fundamental property of self-organizing systems because higher complexity means increasing self-organization.
- Class 1 and 2 cellular automata limiting sets = Regular languages
- Class 3 and 4 cellular automata limiting sets \neq Regular languages.
 - Class 3 cellular automata = context-sensitive languages.
 - Class 4 cellular automata = unrestricted languages(TM).

Computation theory

- Cellular automaton evolution can be viewed as computation.
- Dynamic systems theory can be used to define class 1, 2, and 3 cellular automata; but not for class 4.
- Computational theory needed to define class 4 cellular automata.

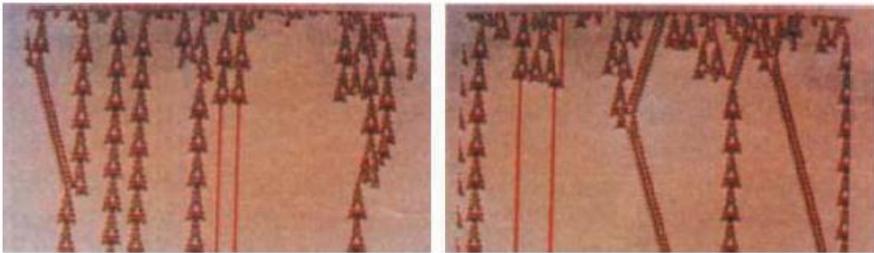


Fig. 5 Examples of the evolution of a typical class 4 cellular automaton from disordered initial states. This and other class 4 cellular automata are conjectured to be capable of arbitrary information processing, or universal computation. The rule shown has $k=3$, $r=1$, and takes the value of a site to be 1 if the sum of the values of the sites in its three-site neighbourhood is 2 or 6, to be 2 if the sum is 3, and to zero otherwise (totalistic code 792).

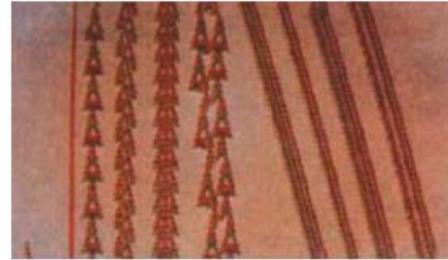


Fig. 6 Persistent structures generated in the evolution of the class 4 cellular automaton of Fig. 5. The first four structures shown have periods 1, 20, 16 and 12 respectively; the last four structures (and their reflections) propagate: the first has period 32, the next three period 3, and the last period 6. These structures are some of the elements required to support universal computation.

- Class 4 cellular automata rules can be used to implement arbitrary information processing operations.
 - Capable of universal computation
 - Evolution can implement any finite algorithm for a given initial conditions

Computation theory

- A short-cut needed for computing the outcome of cellular automaton evolution:
 - Class 1 and 2 need simple computation for future prediction
 - Class 3 and 4 may not allow short-cuts
 - Explicit observation or simulation needed for future prediction.
 - For class 4, simulation is needed for almost all initial configurations. For class 3, some initial configurations might not need simulation.
 - Infinite time limiting behavior will be undecidable (halting problem).
 - Large time limit of entropy can't be computed.

Computation theory

- For class 3 and 4:
 - the occurrence of a particular length n site value sequence in infinite time is undecidable.
 - Decidable in finite time considering initial sequence length $n_0 = n + 2rt$
 - Large n or t will make computation difficult.
- Finding possible sequences generated by class 3 and 4 is an NP-complete problem.
 - Both NP and NP-hard
 - Can't be solved in any time polynomial in n or t .
- Undecidability and intractability are common problems in almost all non-linear systems.
 - No simple formulae for the behavior of such systems (e.g. natural systems)
 - Only simulation and observation are effective for future prediction.

Thermodynamics

- Entropy in conventional thermodynamics (reversible evolution):
 - Fine-grained:
 - Remains constant with time
 - Coarse-grained:
 - Always non-decreasing with time
 - Very little impact on the measures of self-organization structures in long time
- Set of configurations with low coarse-grained entropy
 - Value of every 5th site highly constrained.
 - Arbitrary value for the intervening sites allowed

$$d_{\mu}^{(x)} = - \lim_{X \rightarrow \infty} \frac{1}{X} \sum_{j=1}^{k^x} p_j \log_k p_j$$

Thermodynamics

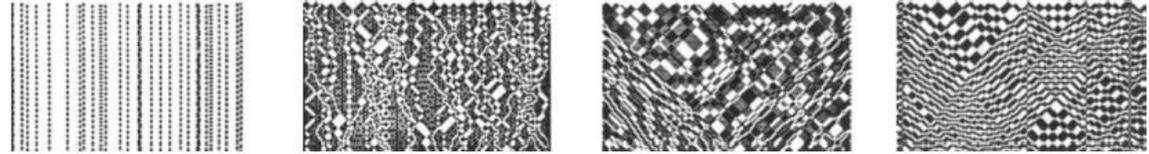


Fig. 7 Evolution of some cellular automata with reversible rules. Each configuration is a unique function of the two previous configurations. (Rule numbers 4, 22, 90 and 126 are shown.) As initial conditions, each site in two successive configurations is chosen to have value 1 with probability 0.1.

- Cellular Automata Rules:

- Reversible:

- Conventional Thermodynamics can be applied
 - Every configuration has a unique predecessor in evolution
 - Spatial entropy and dimension remain constant with time

- Irreversible:

- Spatial entropy and dimension decrease with time
 - Coarse-grained entropy typically increases for a short time, but then decreases to follow the fine-grained entropy

References

Wolfram, S., 1984. Cellular automata as models of complexity. *Nature*, 311(5985), pp.419-424.

Wolfram, S. In "Theory and Applications of Cellular Automata (Including Selected Papers 1983-1986)" [Wolfram, S(Ed.)]. Advanced Series on Complex Systems 1. World Scientific Publishing, 485-557, 1986.

Questions?

Thank you

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